

# On the Construction of Lightweight Circulant Involutory MDS matrices

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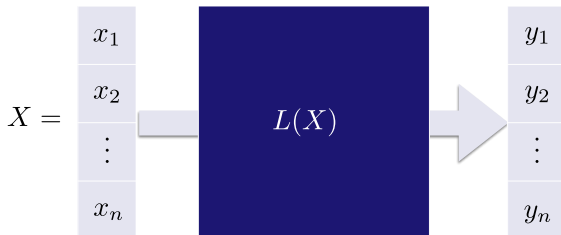
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## Liner Diffusion Layer

A linear mapping provides internal dependency.



The performance is measured by branch number:

- $\mathcal{B}(L) = \min\{\omega_b(X) + \omega_b(L(X)) \mid X \in (\mathbb{F}_2^m)^n, X \neq 0\}$ ;
- $\mathcal{B}(L) \leq n + 1$ .

## Lightweight MDS Matrix

### MDS matrix

Let  $L$  be a linear mapping over  $(\mathbb{F}_2^m)^n$ , i.e.  $L$  is a linear mapping acts on  $n$  S-boxes, and each S-box is of  $m$ -bit length. Then  $L$  is called an MDS matrix if  $\mathcal{B}(L) = n + 1$ .

An MDS matrix can be represented by the following matrix:

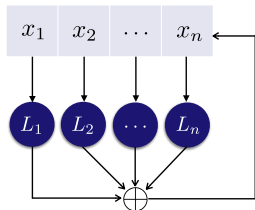
$$L = \begin{pmatrix} L_{1,1} & L_{1,2} & \cdots & L_{1,n} \\ L_{2,1} & L_{2,2} & \cdots & L_{2,n} \\ \vdots & \vdots & \cdots & \vdots \\ L_{n,1} & L_{n,2} & \cdots & L_{n,n} \end{pmatrix},$$

where  $L_{i,j}$  are  $m \times m$  matrices over  $\mathbb{F}_2$ .

**Lightweight** means the implementation requires fewer XORs.

## Previous Construction Methods

- Recursive constructions—Based on LFSRs



- Light but high latency
- Direct constructions—Based on finite fields
  - Implement the multiplication of entries of finite fields

*The entries are always pairwise commutative. Why?*

## The Advantage of Commutative Entries

It is very efficient to characterize whether  $L$  is MDS or not.

Blaum, Roth, IEEE TIT 1999

$L$  is MDS if and only if every square sub-matrices of  $L$  are non-singular.

- The determinants of square sub-matrices of  $L$  can be computed;
- The determinants can be factorized;
- One can get an equivalent condition set.

$$L = \begin{pmatrix} L_{1,1} & L_{1,2} & \cdots & L_{1,n} \\ L_{2,1} & L_{2,2} & \cdots & L_{2,n} \\ \vdots & \vdots & \cdots & \vdots \\ L_{n,1} & L_{n,2} & \cdots & L_{n,n} \end{pmatrix}$$

## The Disadvantages of Commutative Entries

- Some interesting MDS matrices are proved do not exist.
  - Nakahara, Abraho, IJNS 2009:  
 $4 \times 4$  circulant involutory MDS matrix does not exist.
  - Gupta, Ray, CCDS 2015:  
 $n \times n$  circulant involutory MDS matrix does do not exist.
  - Gupta, Ray, CCDS 2015:  
 $2^d \times 2^d$  circulant orthogonal MDS matrix does not exist.

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 $2^d \times 2^d$  circulant orthogonal MDS matrix does not exist over finite fields.
- May lose MDS matrices with fewer XORs.

## Motivations and Methods

*Construct MDS matrices with non-commutative entries.*

### Goals

- Construct some MDS matrices that have been proved do not exist when the entries are commutative.
- Construct MDS matrices with as few XORs as possible.
  - Give lower bounds of XORs of some MDS matrices and corresponding constructions.

### Methods

- Compute the rank of all square sub-matrices to determine whether a matrix is MDS or not.
- Search all possible entries to get lower bounds on XORs.

## Focus on the Following Special Matrices

$$n = 4, m = 4, 8.$$

- Circulant MDS matrix.

- Involutory.
- Non-involutory.
- Orthogonal.

<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>
<i>D</i>	<i>A</i>	<i>B</i>	<i>C</i>
<i>C</i>	<i>D</i>	<i>A</i>	<i>B</i>
<i>B</i>	<i>C</i>	<i>D</i>	<i>A</i>

- Hadamard MDS matrix.

- Involutory.
- Non-involutory.

<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>
<i>B</i>	<i>A</i>	<i>D</i>	<i>C</i>
<i>C</i>	<i>D</i>	<i>A</i>	<i>B</i>
<i>D</i>	<i>C</i>	<i>B</i>	<i>A</i>

- “Optimal”  $4 \times 4$  MDS matrix.

<i>A</i>	<i>I</i>	<i>I</i>	<i>I</i>
<i>I</i>	<i>I</i>	<i>B</i>	<i>A</i>
<i>I</i>	<i>A</i>	<i>I</i>	<i>B</i>
<i>I</i>	<i>B</i>	<i>A</i>	<i>I</i>

## A General Lower Bound on XORs of Entries

The number of entries with no XORs is limited.

### A general bound

Let  $A, B, C, D \in GL(m, \mathbb{F}_2)$ . Then

- If  $Circ(A, B, C, D)$  is a *circulant MDS matrix*, then  $\#A + \#B + \#C + \#D \geq 2$ .
- If  $Had(A, B, C, D)$  is a *Hadamard MDS matrix*, then  $\#A + \#B + \#C + \#D \geq 3$ .

- There are at most 2 entries with no XORs in one row of a circulant MDS matrix;
- There are at most 1 entry with no XORs in one row of a Hadamard MDS matrix.
- We suppose  $L[1, 1] = I$  in the following constructions.

## Trick for Searching

*Using sub-matrices of order 2 to reduce searching space.*

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$$L = \text{Circ}(I, I, A, B) = \begin{pmatrix} I & I & A & B \\ B & I & I & A \\ A & B & I & I \\ I & A & B & I \end{pmatrix}.$$

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$$L = \text{Circ}(I, I, A, B) = \begin{pmatrix} I & I & A & B \\ B & I & I & A \\ A & B & I & I \\ I & A & B & I \end{pmatrix}.$$

### Fact

If  $L = \text{Circ}(I, I, A, B)$  is MDS, then all its square sub-matrices of order 2 are non-singular.

## Trick for Searching

*Using sub-matrices of order 2 to reduce searching space.*

$$\begin{pmatrix} B & A \\ I & I \end{pmatrix}, \begin{pmatrix} A & B \\ B & I \end{pmatrix}, \begin{pmatrix} B & I \\ A & I \end{pmatrix}, \begin{pmatrix} B & A \\ A & I \end{pmatrix}, \\ \begin{pmatrix} I & B \\ B & A \end{pmatrix}, \begin{pmatrix} B & I \\ I & B \end{pmatrix}, \begin{pmatrix} B & I \\ A & B \end{pmatrix}, \begin{pmatrix} A & I \\ I & B \end{pmatrix}, \\ \begin{pmatrix} B & I \\ I & A \end{pmatrix}, \begin{pmatrix} I & A \\ I & I \end{pmatrix}, \begin{pmatrix} I & A \\ A & I \end{pmatrix}, \begin{pmatrix} I & B \\ I & I \end{pmatrix}, \\ \begin{pmatrix} I & B \\ A & I \end{pmatrix}, \begin{pmatrix} A & I \\ I & A \end{pmatrix}, \begin{pmatrix} I & A \\ I & B \end{pmatrix}, \begin{pmatrix} I & I \\ A & B \end{pmatrix}, \\ \begin{pmatrix} I & I \\ I & A \end{pmatrix}, \begin{pmatrix} I & A \\ A & B \end{pmatrix}, \begin{pmatrix} A & B \\ I & I \end{pmatrix}, \begin{pmatrix} I & B \\ I & A \end{pmatrix}, \\ \begin{pmatrix} I & I \\ B & I \end{pmatrix}, \begin{pmatrix} A & B \\ I & A \end{pmatrix}, \begin{pmatrix} I & A \\ B & I \end{pmatrix}, \begin{pmatrix} I & B \\ B & I \end{pmatrix}.$$



## Trick for Searching

*Using sub-matrices of order 2 to reduce searching space.*

### Lemma

*Let  $A, B, C \in GL(m, \mathbb{F}_2)$  are  $m \times m$  non-singular matrices over  $\mathbb{F}_2$ . Then the following statements hold.*

- ①  $\begin{pmatrix} I, & A \\ B, & C \end{pmatrix}$  is non-singular  $\iff \text{rank}(BA + C) = m$ .
- ②  $\begin{pmatrix} A, & I \\ B, & C \end{pmatrix}$  is non-singular  $\iff \text{rank}(CA + B) = m$ .
- ③  $\begin{pmatrix} A, & B \\ I, & C \end{pmatrix}$  is non-singular  $\iff \text{rank}(AC + B) = m$ .
- ④  $\begin{pmatrix} A, & B \\ C, & I \end{pmatrix}$  is non-singular  $\iff \text{rank}(BC + A) = m$ .

## Trick for Searching

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### Conditions

If  $L = \text{Circ}(I, I, A, B)$  is MDS, then the following matrix are non-singular:

$$A + I, B + I, A + B, AB + I, A^2 + B, A + B^2.$$

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$$A, B \in S_{A,B} := \{X \mid \text{rank}(X + I) = m\}.$$

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### Searching method

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for  $A \in S_{A,B}$

$$S_B := \{B \in S_{A,B} \mid \text{rank}(A + B) = \text{rank}(AB + I) \\ = \text{rank}(A^2 + B) = \text{rank}(A + B^2) = m\}$$

for  $B \in S_B$

test whether  $\text{Circ}(I, I, A, B)$  is MDS

---

## Lightweight Circulant Non-involutory MDS Matrix

Searching: Magma v2.20-3, Laptop (OSX 10.9, Intel Core i7);

It takes nearly *3 days* to get the following result:

### Theorem

Let  $L = \text{Circ}(I, I, A, B)$ , where  $A, B \in GL(m, \mathbb{F}_2)$ ,  $m = 4, 8$ . If  $L$  is an MDS matrix, then  $\#A + \#B \geq 3$ . (48, 80640)

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(48, 80640) means there are at least 48 and 80640 pairs of  $A, B \in GL(m, \mathbb{F}_2)$  such that  $\text{Circ}(I, I, A, B)$  is MDS and the lower bound on XORs holds for  $m = 4$  and  $m = 8$  respectively.

## Lightweight Circulant Involutory MDS Matrices

### Equivalent conditions of circulant involutory matrix

Let  $L = \text{Circ}(I, A, B, C)$ , where  $A, B, C \in GL(m, \mathbb{F}_2)$ . Then  $L$  is an involution if and only if:

$$AB = BA, BC = CB, A^2 = C^2, AC + CA = B^2.$$



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### A general construction

- $A, C \in GL(m, \mathbb{F}_2)$  with  $A^2 = C^2 = I$
- The multiplication order of  $A + C$  equals  $4k - 2$ , where  $k > 1$ . This means  $(A + C)^{4k-2} = I$ .
- Let  $B = (A + C)^{2k}$

Then the matrix  $\text{Circ}(I, A, B, C)$  is an involution.

## Lightweight Circulant Involutory MDS Matrices

Searching: Magma v2.20-3, Laptop (OSX 10.9, Intel Core i7);

It takes nearly *one week* to get the following result:

### Theorem

*Their exist circulant involutory MDS matrix over  $GL(m, \mathbb{F}_2)$ ,  $m = 4, 8$ . Furthermore, the following statements hold.*

- ①  $m = 4$ , circulant involutory MDS matrix  $Circ(I, A, B, C)$  satisfies  $\#A + \#B + \#C \geq 5$ . (48)
- ②  $m = 8$ , if  $\#A \leq 3$  and  $\#C \leq 3$ , then circulant involutory MDS matrix  $Circ(I, A, B, C)$  constructed with the above method satisfies  $\#A + \#B + \#C \geq 9$ . (40320)
- ③  $m = 4, 8$ , circulant involutory MDS matrix  $Circ(I, A, B, B, A)$  satisfies  $\#A + \#B \geq \frac{m}{2}$ . (24, 20160)

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- Circulant orthogonal MDS matrix exists over  $GL(m, \mathbb{F}_2)$ .

## Lightweight Hadamard Involutory MDS Matrices

### Equivalent conditions of Hadamard involutory matrix

Let  $A, B, C \in GL(m, \mathbb{F}_2)$ . Then  $L = Had(I, A, B, C)$  is an involution if and only if  $A, B, C$  are pairwise commutative and  $A^2 + B^2 = C^2$ .

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Searching: Magma v2.20-3, Laptop (OSX 10.9, Intel Core i7);  
It takes about *one day* to get the following result:

### Theorem

- 1 Let  $A, B, C \in GL(4, \mathbb{F}_2)$ . If  $L = Had(I, A, B, C)$  is an MDS involutory matrix, then  $\#A + \#B + \#C \geq 6$ . (144)
- 2 Let  $A \in GL(8, \mathbb{F}_2)$  with  $\#A \leq 3$ . If  $L = Had(I, A, A^{-1}, A + A^{-1})$  is an MDS involutory matrix, then  $\#A + \#A^{-1} + \#(A + A^{-1}) \geq 10$ . (80640)

## Lightweight Hadamard Non-involutory MDS matrices

Searching: Magma v2.20-3, PC (Win7, Intel Core i5);

It takes nearly 4 weeks to get the following result:

### Theorem

- 1 Let  $A, B, C \in GL(4, \mathbb{F}_2)$ . If  $L = \text{Had}(I, A, B, C)$  is an MDS matrix, then  $\#A + \#B + \#C \geq 4$ . (72)
- 2 Let  $A, B \in GL(8, \mathbb{F}_2)$ . If  $L = \text{Had}(I, A, A^T, B)$  is an MDS matrix, then  $\#A + \#A^T + \#B \geq 5$ . (622)

## Lightweight “Optimal” $4 \times 4$ MDS Matrix

$$L = \begin{pmatrix} A & I & I & I \\ I & I & B & A \\ I & A & I & B \\ I & B & A & I \end{pmatrix}$$

### Theorem

Let  $L$  be a matrix constructed as above, where  $A, B \in GL(m, \mathbb{F}_2)$ ,  $m = 4, 8$ . If  $L$  is an MDS matrix, then

$$4\#A + 3\#B \geq \begin{cases} 13, & m = 4; \text{ (24)} \\ 10, & m = 8. \text{ (40320)} \end{cases}$$

## Comparisons with Previous Constructions

### Comparisons of non-involutory MDS matrices

Matrix	Entries	The first row	XOR count	Ref.
Circulant	$GL(8, \mathbb{F}_2)$	$(I, I, A, B)$	$3+3 \times 8=27$	Our paper
Circulant	$\mathbb{F}_{2^8}/0x11b$	$(0x02, 0x03, 0x01, 0x01)$	$14+3 \times 8=38$	AES
Hadamard	$GL(8, \mathbb{F}_2)$	$(I, A, A^T, B)$	$5+3 \times 8=29$	Our paper
Hadamard	$\mathbb{F}_{2^8}/0x1c3$	$(0x01, 0x02, 0x04, 0x91)$	$13+3 \times 8=37$	Sim et al.
Subfield-Hadamard	$\mathbb{F}_{2^4}/0x13$	$(0x1, 0x2, 0x8, 0x9)$	$2 \times (5+3 \times 4)=34$	Sim et al.



## Comparisons with Previous Constructions

### Comparisons of involutory MDS matrices

Matrix	Entries	The first row	XOR count	Ref.
Circulant	$GL(8, \mathbb{F}_2)$	$(I, I, A, B)$	$9+3 \times 8=33$	Our paper
Hadamard	$GL(8, \mathbb{F}_2)$	$(I, A, A^{-1}, A + A^{-1})$	$10+3 \times 8=34$	Our paper
Subfield-Hadamard	$\mathbb{F}_{2^4}/0x13$	$(0x1, 0x4, 0x9, 0xd)$	$2 \times (6+3 \times 4)=36$	Sim et al.
Hadamard	$\mathbb{F}_{2^8}/0x165$	$(0x01, 0x02, 0xb0, 0xb2)$	$16+3 \times 8=37$	Sim et al.
Hadamard	$\mathbb{F}_{2^8}/0x11d$	$(0x01, 0x02, 0x04, 0x06)$	$22+3 \times 8=46$	Anubis
Compact Cauchy	$\mathbb{F}_{2^8}/0x11b$	$(0x01, 0x12, 0x04, 0x16)$	$54+3 \times 8=78$	Cui et al.
Hadamard-Cauchy	$\mathbb{F}_{2^8}/0x11b$	$(0x01, 0x02, 0xfc, 0xfe)$	$74+3 \times 8=98$	Gupta et al.

## Comparisons with Previous Constructions

Comparisons of MDS matrices over  $\mathbb{F}_2^4$  and  $\mathbb{F}_{2^4}$

Matrix	Entries	The first row	XOR count	Ref.
Circulant	$GL(4, \mathbb{F}_2)$	$(I, A, B, C)$	<b>3</b> +3×4=15	Our paper
Involutory circulant	$GL(4, \mathbb{F}_2)$	$(I, A, B, C)$	<b>5</b> +3×4=17	Our paper
Hadamard	$GL(4, \mathbb{F}_2)$	$(I, A, B, C)$	<b>4</b> +3×4=16	Our paper
Hadamard	$\mathbb{F}_{2^4}/0x13$	$(0x1, 0x2, 0x8, 0x9)$	5+3×4=17	Sim et al.
Involutory Hadamard	$GL(4, \mathbb{F}_2)$	$(I, A, A^{-1}, A + A^{-1})$	<b>6</b> +3×4=18	Our paper
Involutory Hadamard	$\mathbb{F}_{2^4}/0x13$	$(0x1, 0x4, 0x9, 0xd)$	6+3×4=18	Joltik
Involutory Hadamard	$\mathbb{F}_{2^4}/0x19$	$(0x1, 0x2, 0x6, 0x4)$	6+3×4=18	Prøst

## Summary

*Construct circulant involutory MDS matrices firstly.*

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MDS Matrix	$m$	lower bounds
Involutory $Circ(I, A, B, C)$	4, 8	$\#A + \#B + \#C \geq m + 1$
Involutory $Circ(I, A, B, B, A)$	4, 8	$\#A + \#B \geq m/2$
Involutory $Had(I, A, B, C)$	4, 8	$\#A + \#B + \#C \geq m + 2$
Non-involutory $Circ(I, I, A, B)$	4, 8	$\#A + \#B \geq 3$
Non-involutory $had(I, A, B, C)$	4	$\#A + \#B + \#C \geq 4$
Non-involutory $had(I, A, A^T, B)$	8	$\#A + \#A^T + \#C \geq 5$
Orthogonal $Circ(I, A, B, C)$	4	$\#A + \#B + \#C \geq 8$
“optimal” $4 \times 4$ MDS matrix	4	$4\#A + 3\#B \geq 13$
“optimal” $4 \times 4$ MDS matrix	8	$4\#A + 3\#B \geq 10$

The lower bounds on XORs of “red” MDS matrices can not be improved.

*Is the bound of circulant and Hadamard MDS matrices can be improved if we remove the supposition  $L[1, 1] = I$ ?*

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The answer for MDS matrix over  $GL(4, \mathbb{F}_2)$  is NO.

## Lower bounds on MDS matrix over $GL(4, \mathbb{F}_2)$

Searching: Magma v2.20-3, PC (Win7, Intel Core i5).

After about *50 days* searching, we get the following result.

### Theorem

Let  $A_1, A_2, A_3, A_4 \in GL(4, \mathbb{F}_2)$ , and  $\mathcal{B} = \sum_{i=1}^4 \#A_i$ .

- 1 If  $\text{Circ}(A_1, A_2, A_3, A_4)$  is an MDS matrix, then  $\mathcal{B} \geq 3$ ;
- 2 If  $\text{Circ}(A_1, A_2, A_3, A_4)$  is an involutory MDS matrix, then  $\mathcal{B} \geq 5$ ;
- 3 If  $\text{Had}(A_1, A_2, A_3, A_4)$  is an MDS matrix, then  $\mathcal{B} \geq 4$ ;
- 4 If  $\text{Had}(A_1, A_2, A_3, A_4)$  is an involutory MDS matrix, then  $\mathcal{B} \geq 6$ ;

## Future Problems

- ① Construct bigger order, such as  $8 \times 8$ , lightweight MDS matrices with non-commutative entries.
- ② Characterize the existence of  $8 \times 8$  circulant MDS matrix over  $GL(4, \mathbb{F}_2)$ .
  - *It has been verified that  $8 \times 8$  circulant MDS matrix do not exist over  $\mathbb{F}_{2^4}$  [Khoo et al. CHES 2014].*



Thank you!